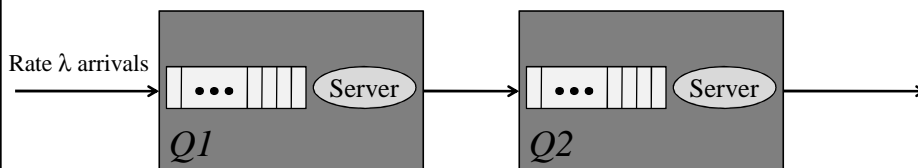


Return to Queueing Theory

- Build on the previous material
 - From single-hop queueing delay to end-to-end queueing delay
 - Overall delay experienced in routing a packet from source to destination node
- Queueing Networks
 - Section 3.6 in text
 - Each queue represents a node which may have multiple input streams of “arrivals” from other nodes
 - Also, output stream may be split to multiple next-hop queues
 - Arrival process at one queue may depend on departure processes of multiple other queues, so probably is not Poisson
 - System is complex! What can be done?

1

Example

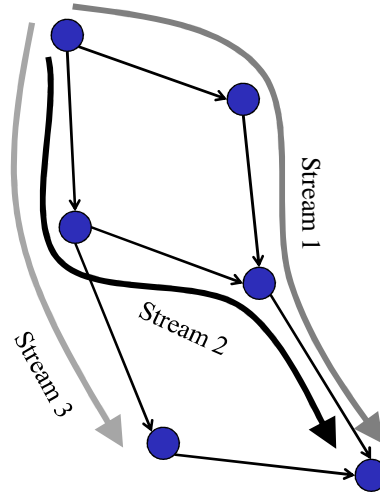


- “Transmission line” model
 - Rate λ arrivals over one link, fixed length packets
 - Service times are constant (equal to transmission delay $1/\mu$), so $Q1$ is $M/D/1$ (D means deterministic) Arrivals into $Q2$ separated by at least $1/\mu$
 - If rate of second link is \geq that of first link, there is ***NO waiting in the second queue EVER!***
 - If packet lengths are iid exponential, this becomes $M/M/1$
 - How to characterize/analyze the second queue?

2

How to Evaluate Delay?

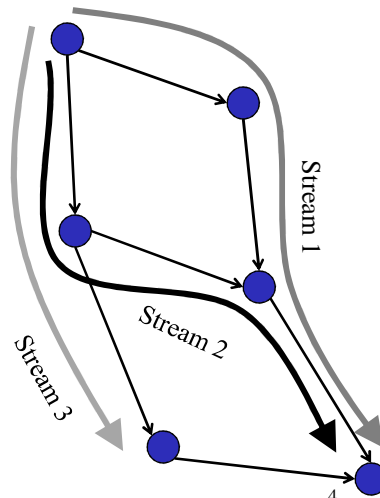
- Setup:
 - Assume multiple packet streams each following a unique path
 - Let x_s be the arrival rate (in packets/second) for stream s
 - For circuit switching:
 - Arrivals on a link (i,j) occur with rate λ_{ij} equal to summation of x_s for each stream s traversing (i,j)
 - For packet switching:
 - Include fraction $f_{ij}(s)$ of packets in stream s traversing (i,j)



3

Kleinrock Independence Approximation

- 1964: Kleinrock suggests that random merger of several packet streams into a single stream restores independence of arrival times and packet lengths
- In particular, for densely connected networks with moderate-to-heavy traffic, inputs are approximately Poisson processes



4

Analysis under KIA

- Average number of packets in queue or service:
 - For link (i,j) , $N_{ij} = \lambda_{ij} / (\mu_{ij} - \lambda_{ij})$ where $1/\mu_{ij}$ is average transmit delay
 - Total packets in queue or service N is summation of N_{ij} over all links (i,j)
- Average queueing delay:
 - From Little's Theorem, $T = N/\gamma$, where γ is sum rate of all x_s arrival streams
- Including proc/prop delay d_{ij} at each (i,j) :
 - Add the term $\lambda_{ij}d_{ij}$ in the summation

5

Time Reversibility

- Let X be a DTMC, and let X_n denote the state at time $n \gg 0$. Let's look at the behavior backward in time, instead of looking forward in time.
- Let P^*_{ij} be the reverse transition probability given by $\Pr[X_m = j \mid X_{m+1} = i]$.
- Show that $\pi_i P^*_{ij} = \pi_j P_{ji}$
- *Def:* A DTMC X is **time reversible** if $P^*_{ij} = P_{ij}$. In this case, the stationary distribution of the reversed DTMC X^* is the same as the forward chain X .
 - We immediately see that a DTMC X is time reversible *if and only if* the detailed balance equations hold.
- This extends in a straightforward way to CTMCs.

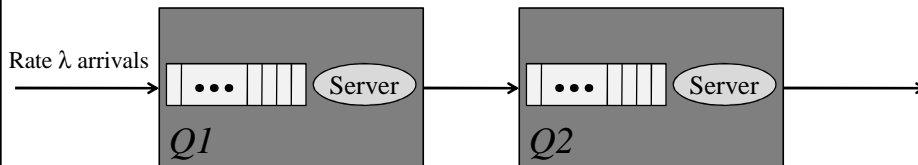
6

Burke's Theorem

- *Theorem (Burke)*: Consider an $M/M/1$, $M/M/k$, or $M/M/\infty$ queue with arrival rate λ , and suppose the system starts in steady-state. Then:
 - The departure process is Poisson with rate λ .
 - The number of customers in the system at time t is independent of the sequence of departure times prior to t .
- Why is this true?
 - $M/M/1$, $M/M/k$, and $M/M/\infty$ are time reversible CTMCs.
 - Departures prior to time t are arrivals after time t in reversed process, which is Poisson, so future arrivals do not depend on system occupancy.

7

Example – Revisited



- Returning to the previous example:
 - Suppose $Q1$ is $M/M/1$ with exp. service rate μ_1 .
 - From Burke's Theorem, arrivals to $Q2$ are Poisson with rate λ and exp. service rate $\mu_2 \rightarrow M/M/1$.
 - $\rightarrow \pi_n^{(i)} = \rho_i^n (1 - \rho_i)$, $\rho_i = \lambda / \mu_i$, $i=1,2$
 - From 2nd part of Burke's Theorem, number of customers in each queue is independent, so $\Pr[n \text{ in } Q1, m \text{ in } Q2] = \pi_n^{(1)} \pi_m^{(2)}$
 - This is a simple example of an *acyclic queueing network*

8

Queueing Networks

- Setup:
 - A network of K queues, each with a single server, service rate μ_i at queue i
 - Arrivals into the network at queue i follow a Poisson process with rate r_i . At least one r_i is positive.
 - Customer served by queue i moves to queue j with probability P_{ij} – leaves with residual probability
 - Total arrival rate into queue j is

$$\lambda_j = r_j + \sum_{i=1}^K \lambda_i P_{ij}, \quad \text{for } j = 1, \dots, K$$
 - Let $\rho_j = \lambda_j / \mu_j$

9

Queueing Network State Space

- Let $n = (n_1, \dots, n_K)$ denote a state in the state space \mathcal{X}^K , where \mathcal{X} is the state space of each queue (here, the non-negative integers), i.e. n_i = number of customers in queue i .
- State transitions (with very high probability):
 - New arrival at queue j (with rate r_j)
 - State $n(j^+) = n + e_j = (n_1, \dots, n_{j-1}, n_j + 1, n_{j+1}, \dots, n_K)$
 - Exit system from queue j (with rate $\mu_j \left(1 - \sum_{i=1}^K P_{ji}\right)$)
 - State $n(j^-) = n - e_j = (n_1, \dots, n_{j-1}, n_j - 1, n_{j+1}, \dots, n_K)$
 - Customer moves from j to i (with rate $\mu_j P_{ji}$)
 - State $n(j^-, i^+) = n - e_j + e_i$

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Jackson's Theorem

- *Theorem (Jackson)*: For such a queueing network, if $\rho_j < 1$ for $j=1, \dots, K$, then for all $n=(n_1, \dots, n_K)$, $n_j \geq 0$, we have

$$\pi_n = \prod_{j=1}^K \pi_{n_j}^{(j)}$$

$$\pi_{n_j}^{(j)} = \rho_j^{n_j} (1 - \rho_j)$$

- In other words, queueing network behaves as **independent collection of M/M/1 queues** even though the arrival to each queue **is not Poisson**.

11

Closed Queueing Networks

- Same setup as before, only $r_i=0$ and P_{ij} sum to 1 over j for all $i=1, \dots, K$, i.e. nobody enters and nobody leaves, so

$$\lambda_j = \sum_{i=1}^K \lambda_i P_{ij}, \quad \text{for } j = 1, \dots, K$$

and total number of customers is fixed at M.

- Here, π_n can be positive only if

$$n_1 + \dots + n_K = M.$$

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